



Supporting nested resources in MrsP

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Sporadic task model

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- Preemptive fixed-priority scheduling

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- Preemptive fixed-priority scheduling
- Follows PCP:
 - Resources have a Local Ceiling Priority
 - A job is (locally) blocked as most once during its execution
 - Before executing
 - Then all resources are (locally) available
 - Deadlocks are prevented

$$R_i = C_i + \max(\hat{e}, \hat{b}) + \sum_{\tau_j \in \mathbf{hpl}(i)} \left\lceil \frac{R_i}{T_j} \right\rceil C_j$$

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$$C_i = WCET_i + \sum_{r^j \in \mathbf{F}(\tau_i)} n_i e^j$$

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 $e^j = |map(G(r^j))|c^j$

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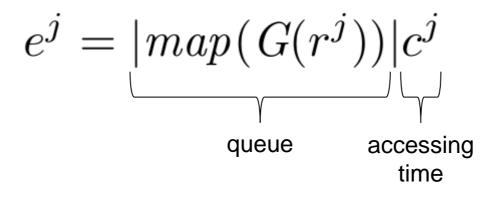
$$e^j = |map(G(r^j))|c^j$$

 $e^{j} = |map(G(r^{j}))|c^{j}$ queue

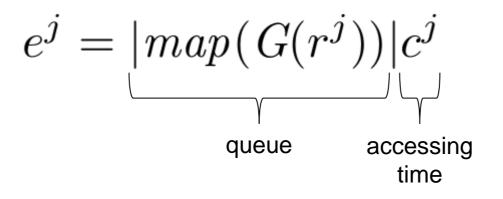
$$e^{j} = |map(G(r^{j}))|c^{j}$$

One access request per processor

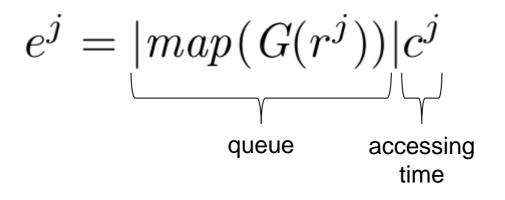
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One access request per processor

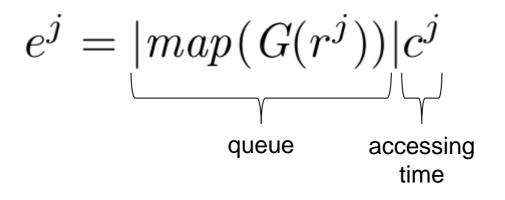


- One access request per processor
- Access time equal to execution time

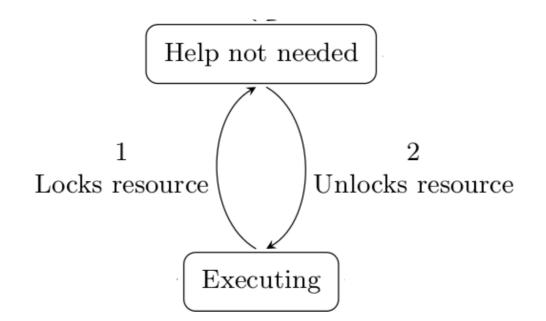


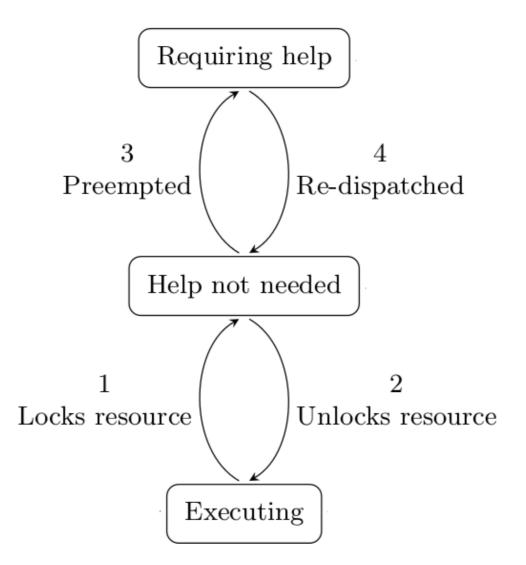
- One access request per processor
- Access time equal to execution time
 - Resource access preemptable

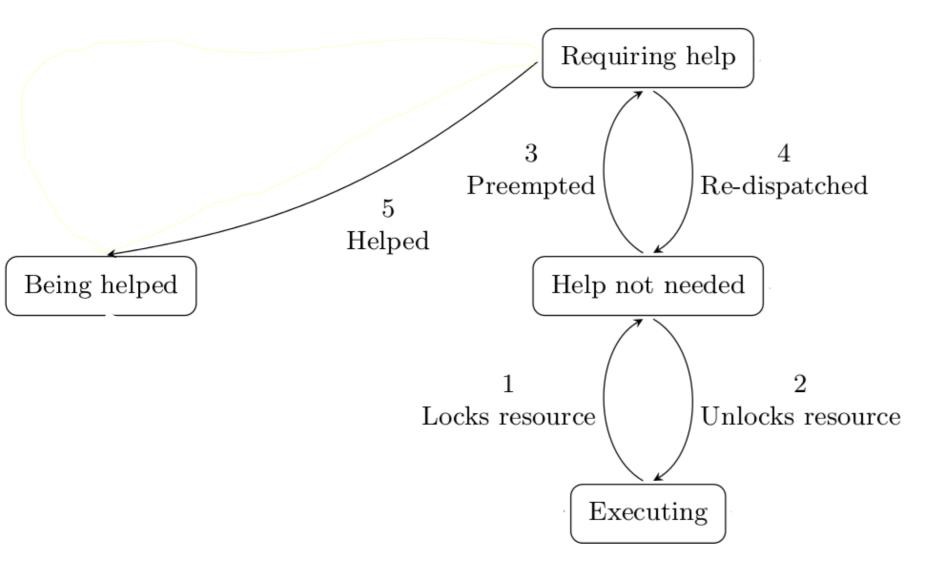
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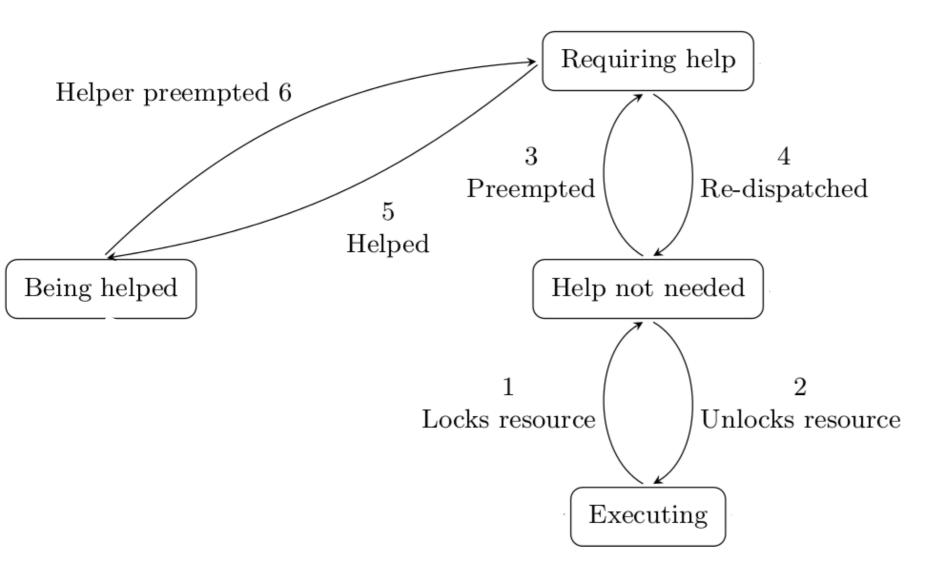


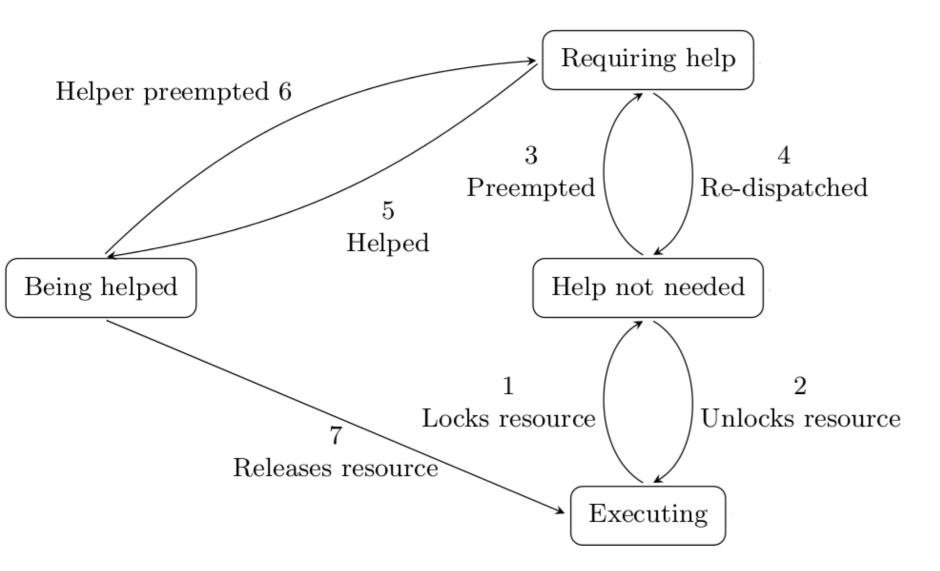
- One access request per processor
- Access time equal to execution time
 - Resource access preemptable
 - Helping mechanism

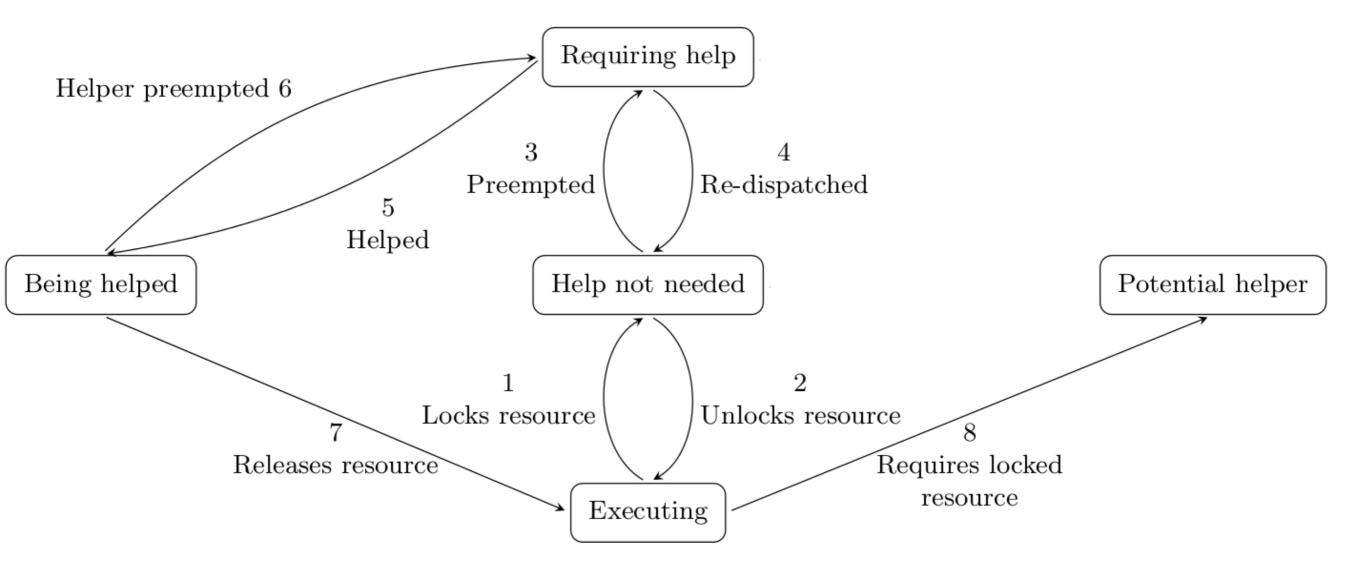


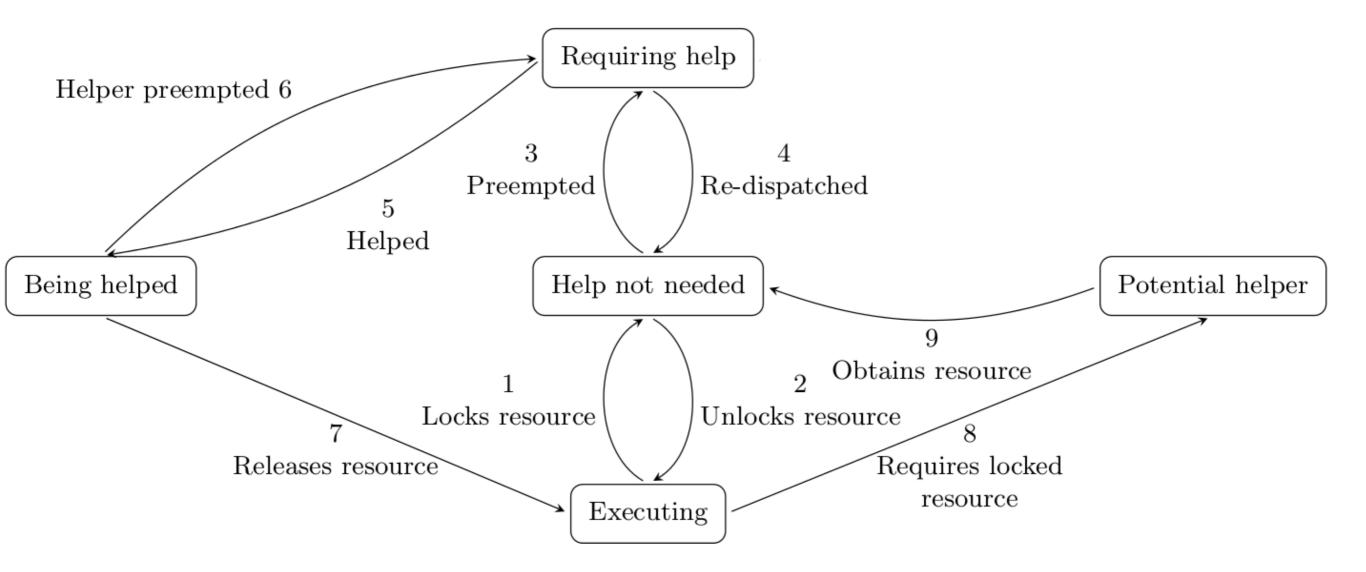


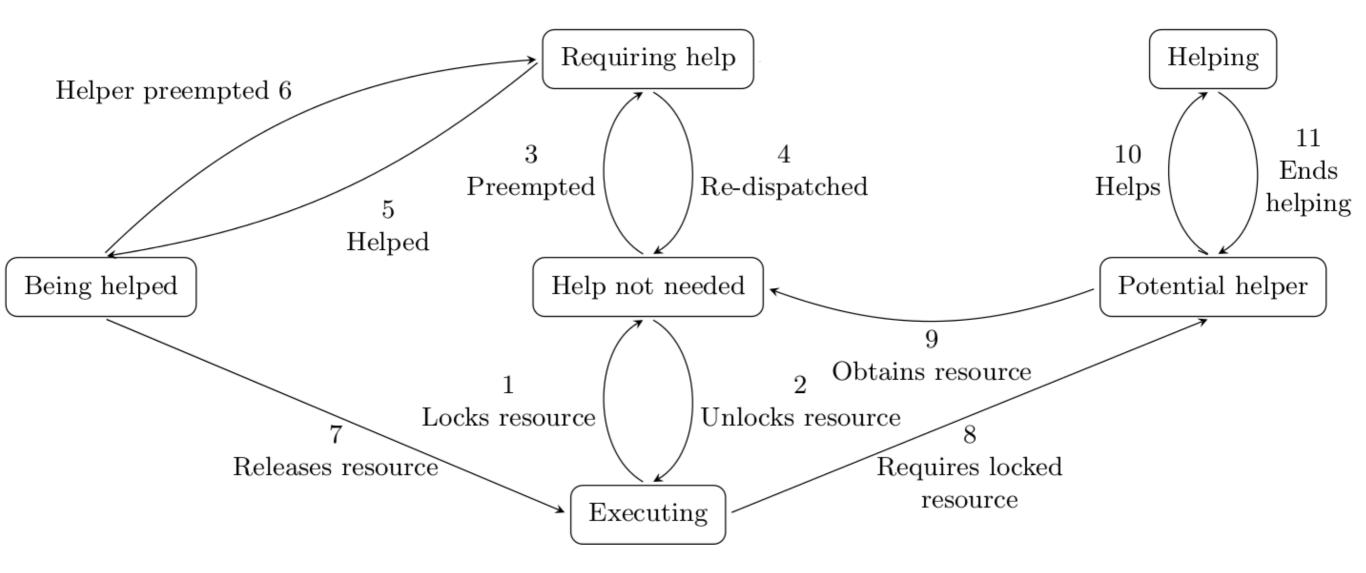


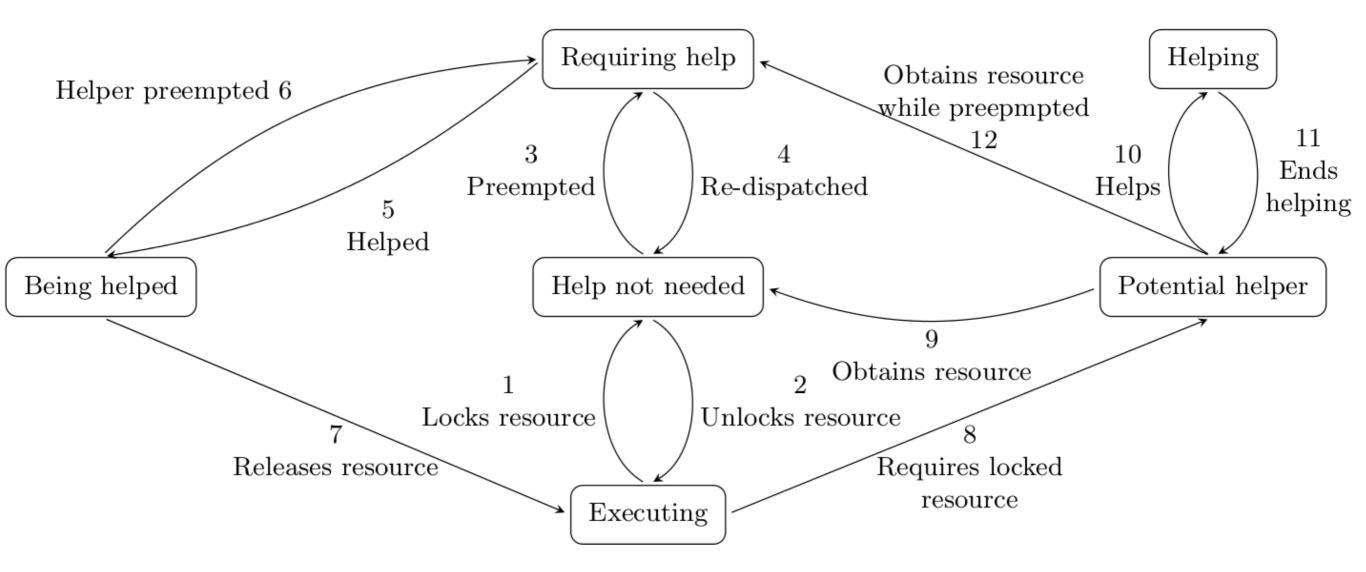


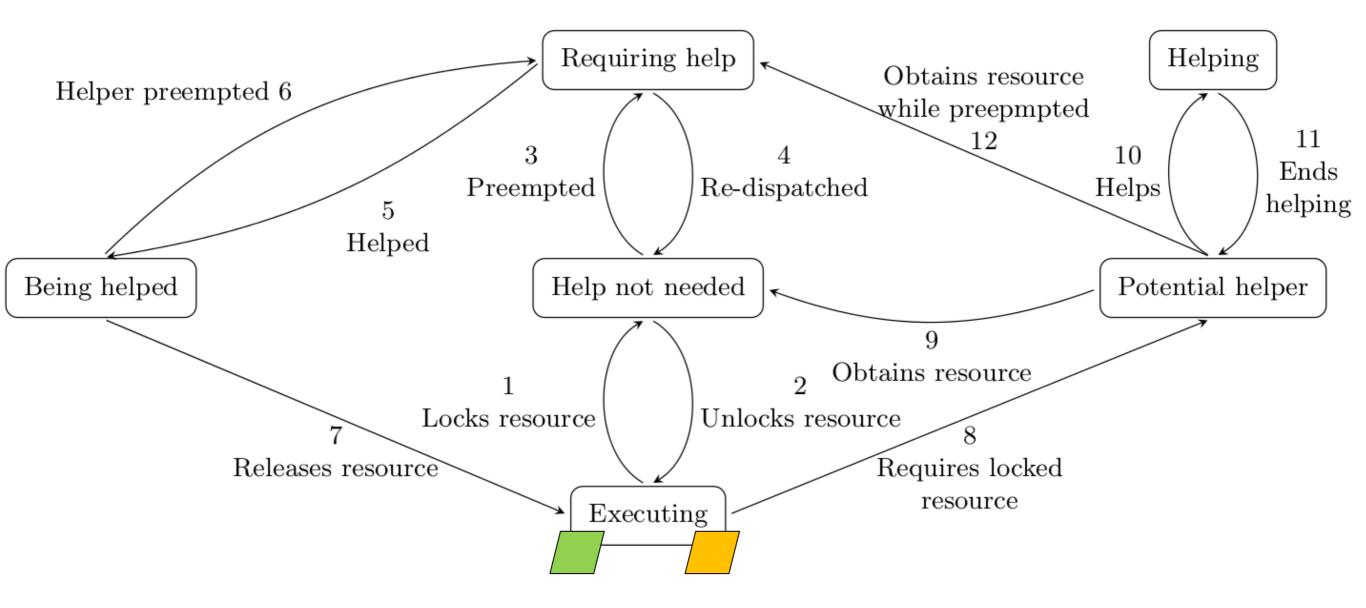


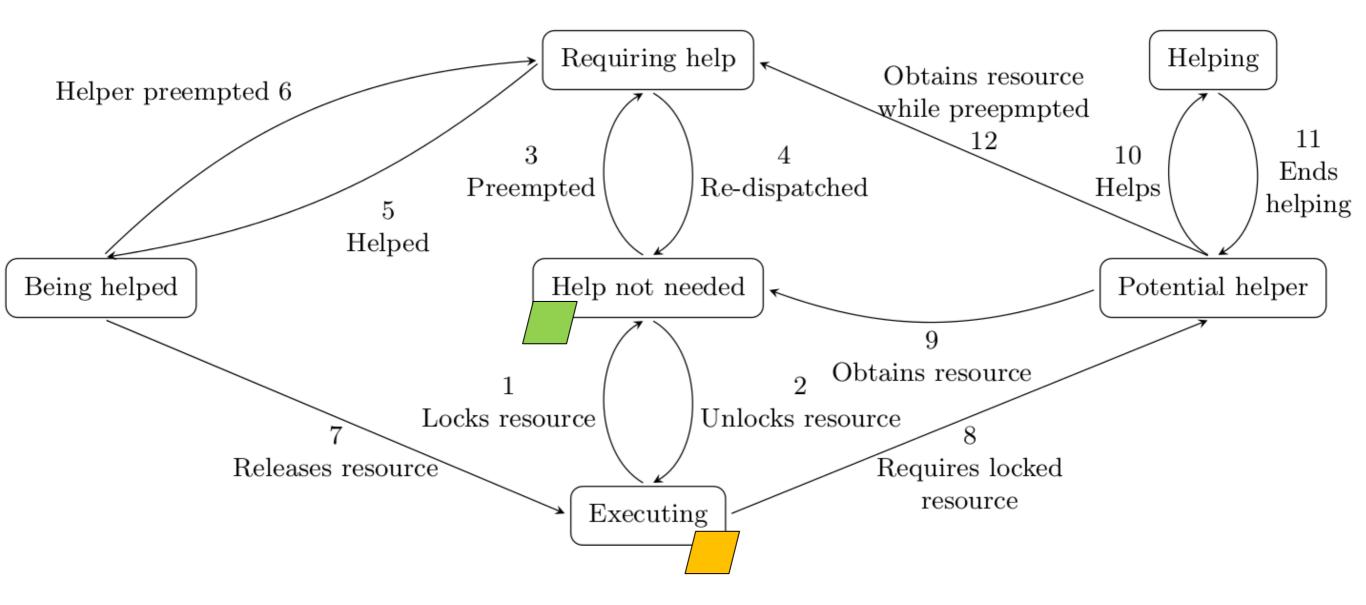


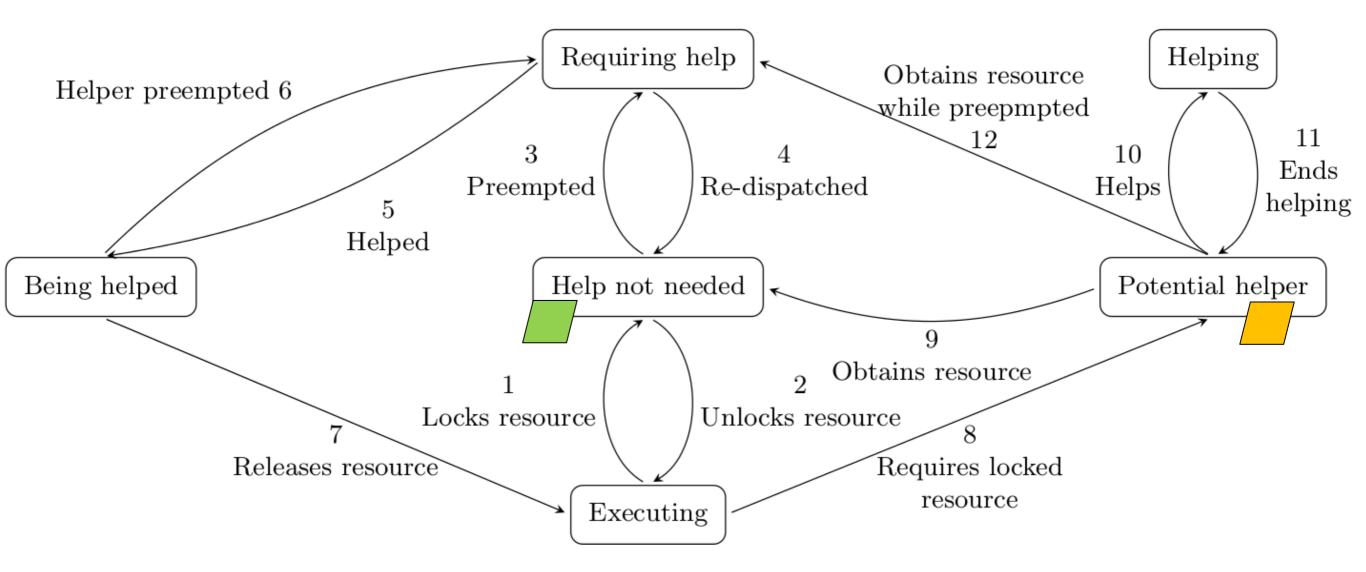


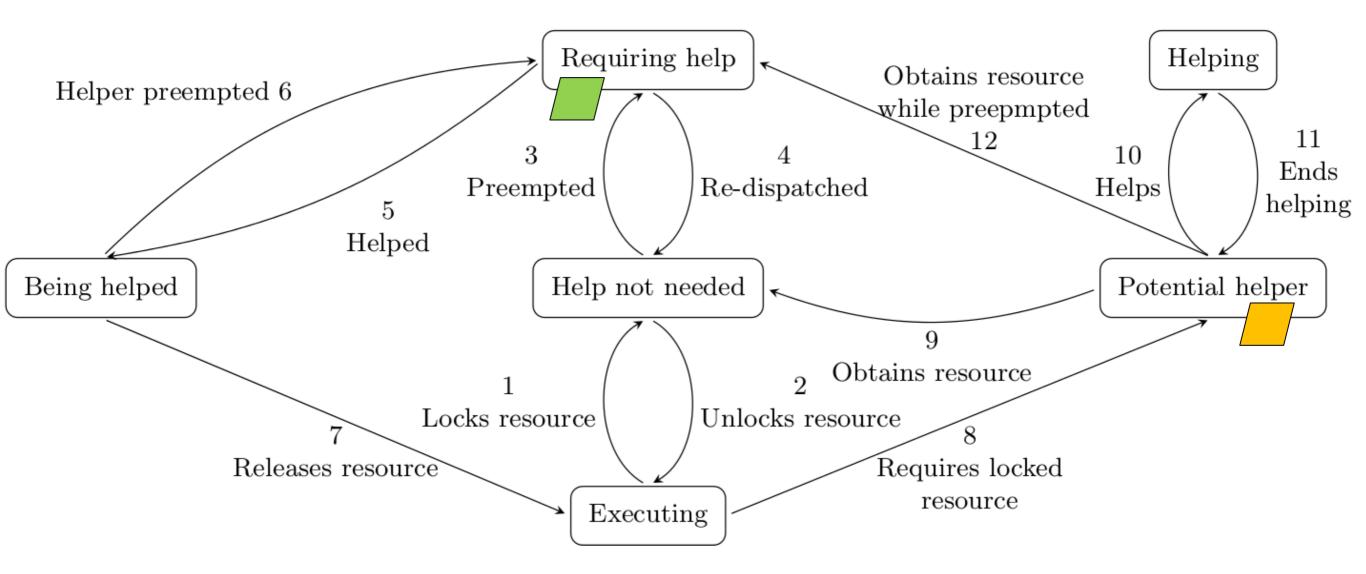


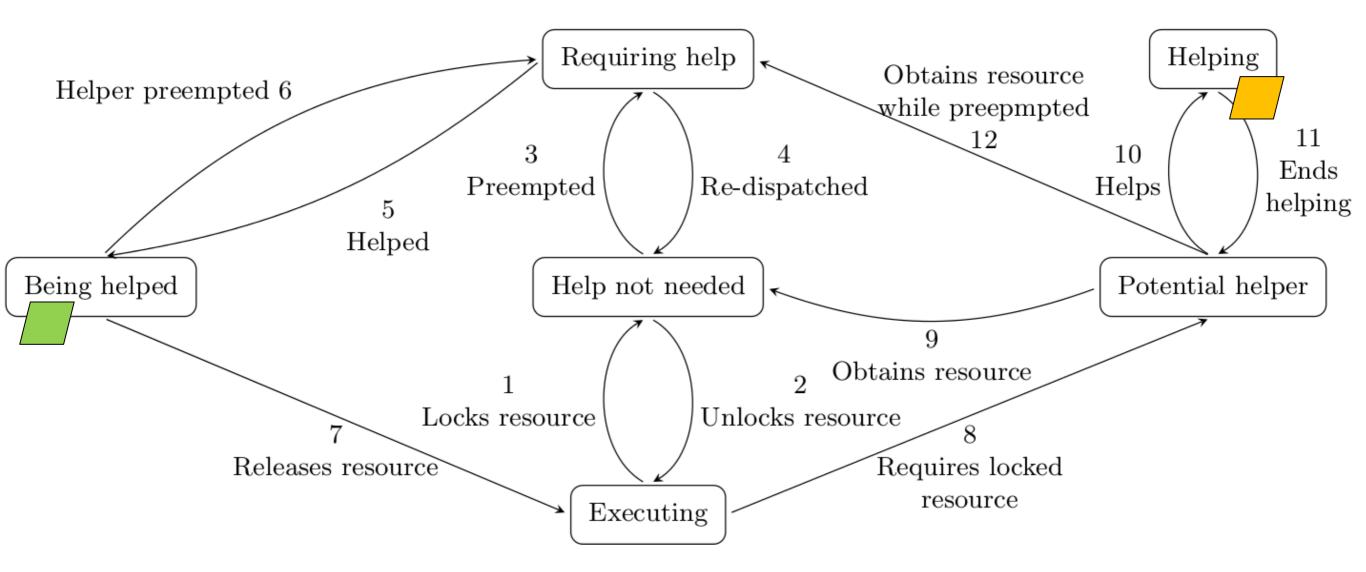


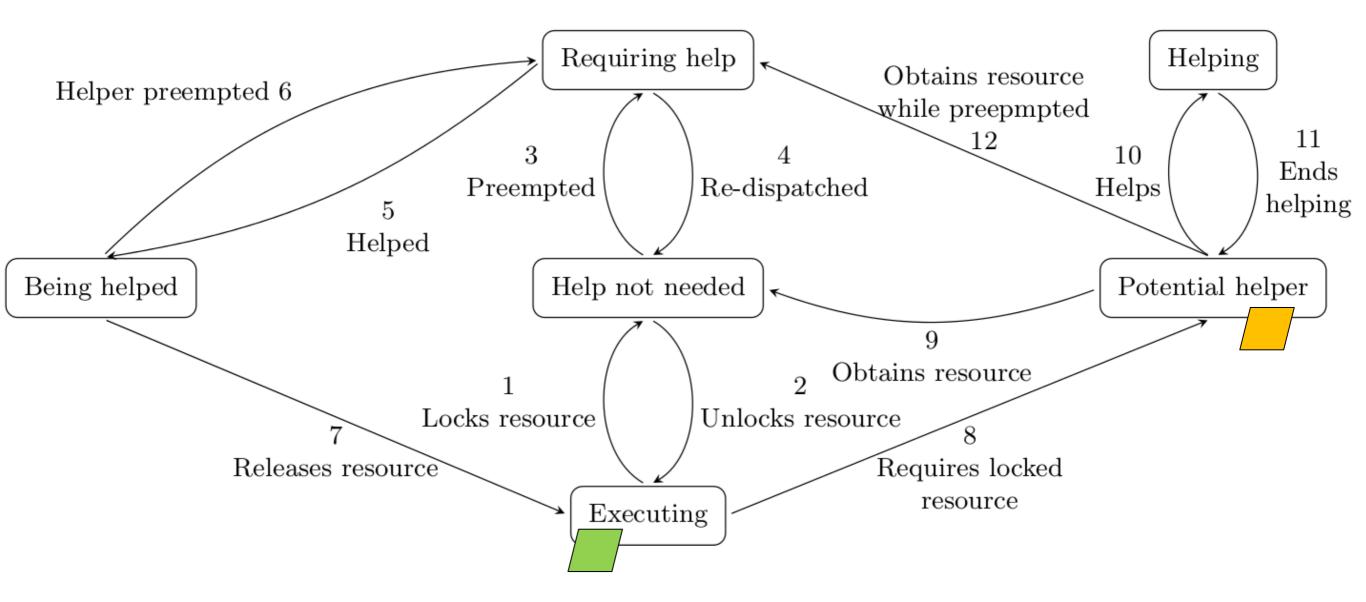


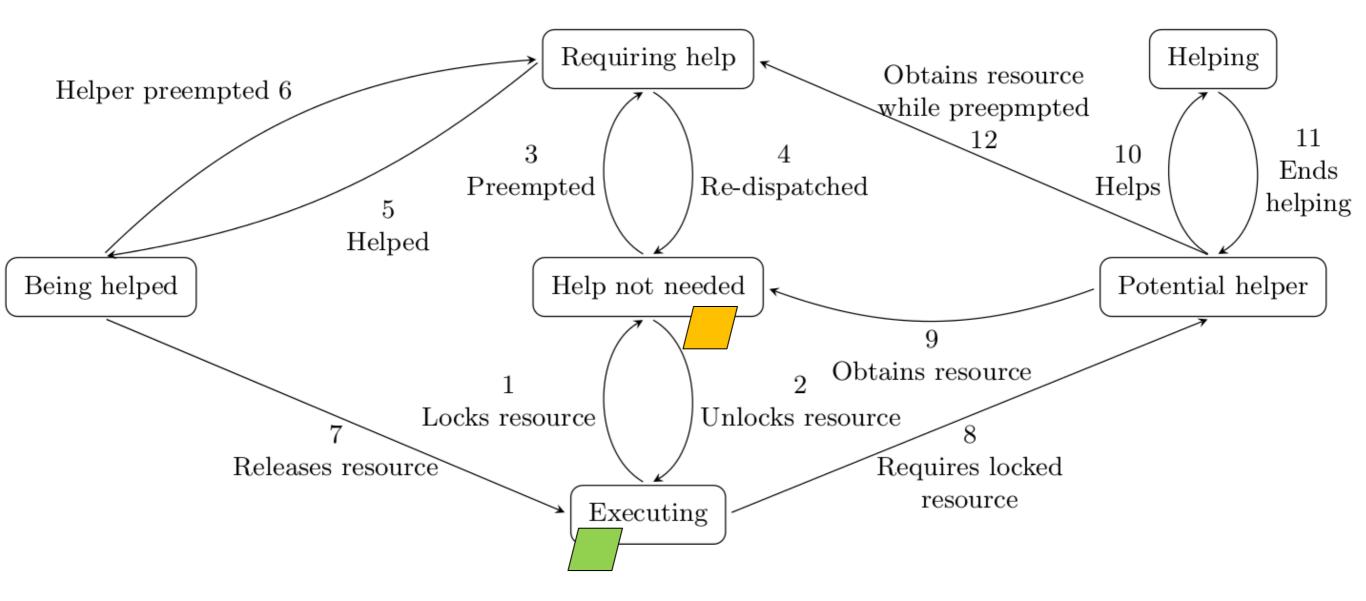












MrsP – Nested Model

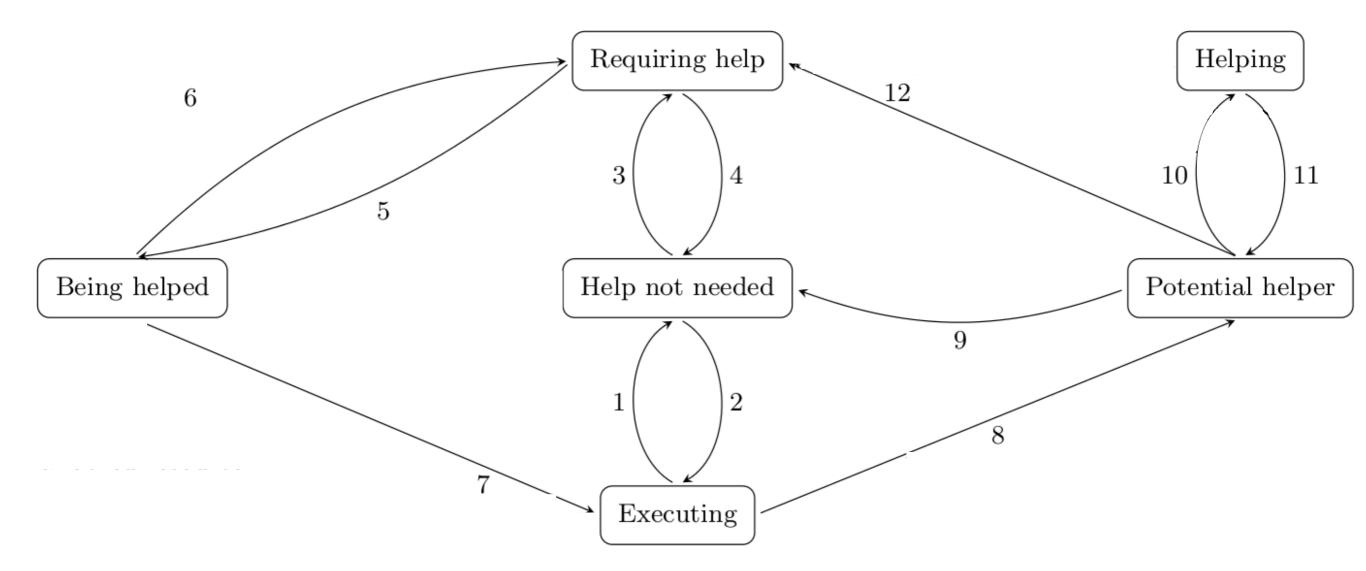
Challenges

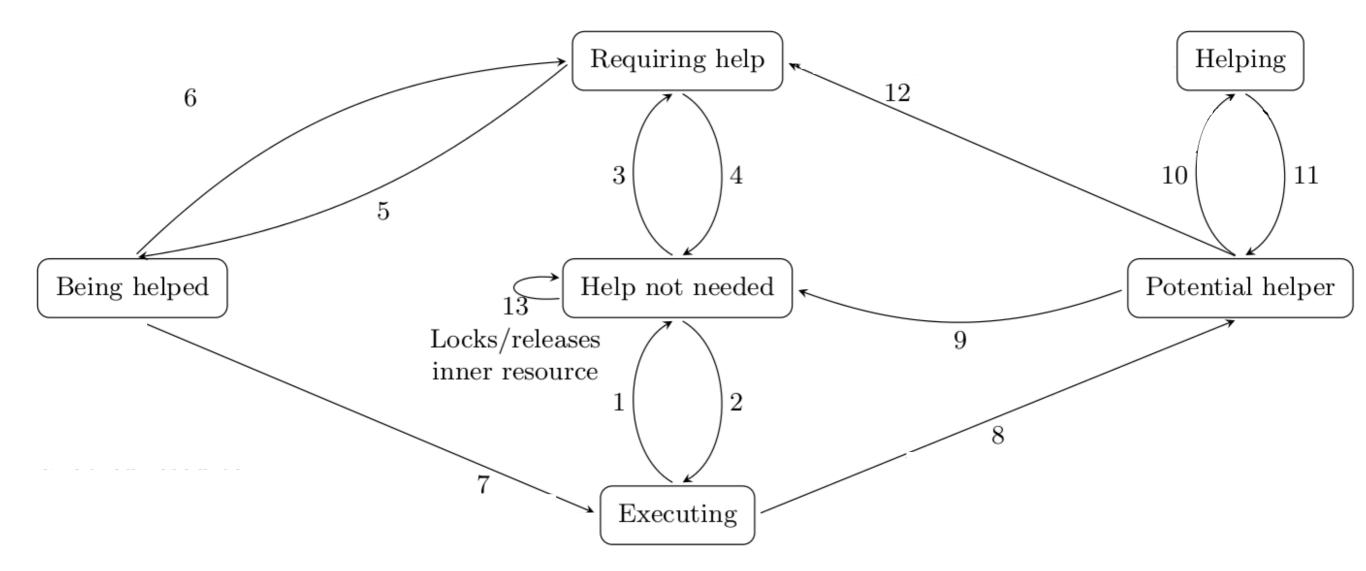
- Challenges
 - Global and Local resources

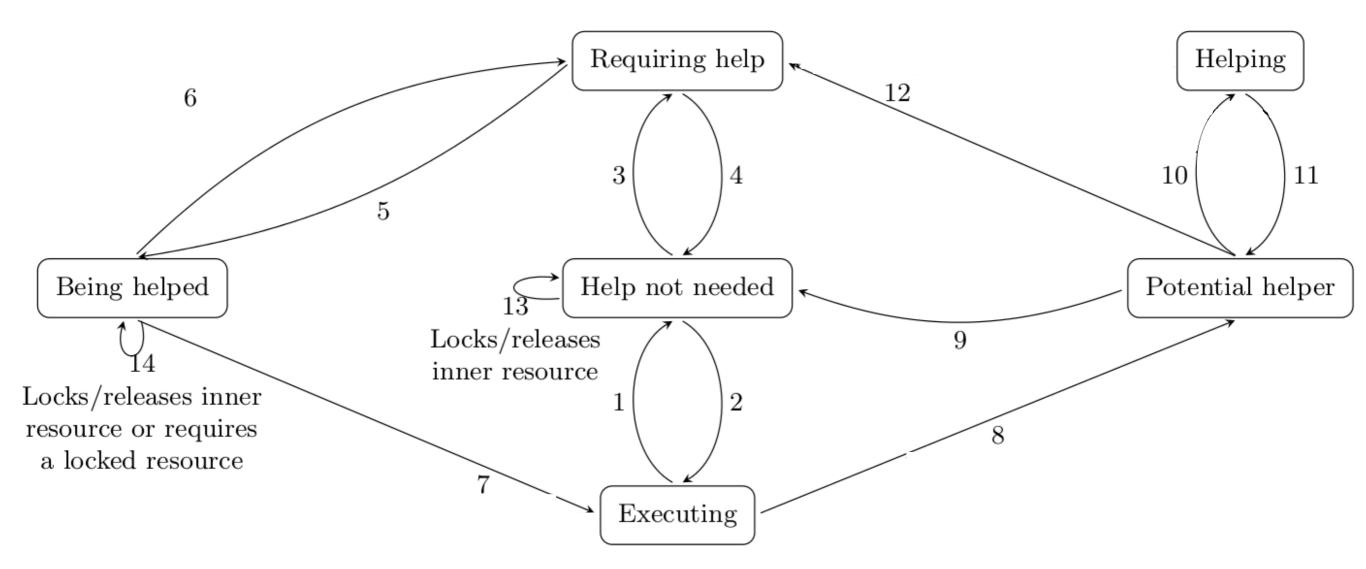
- Challenges
 - Global and Local resources
 - Bound concurrency

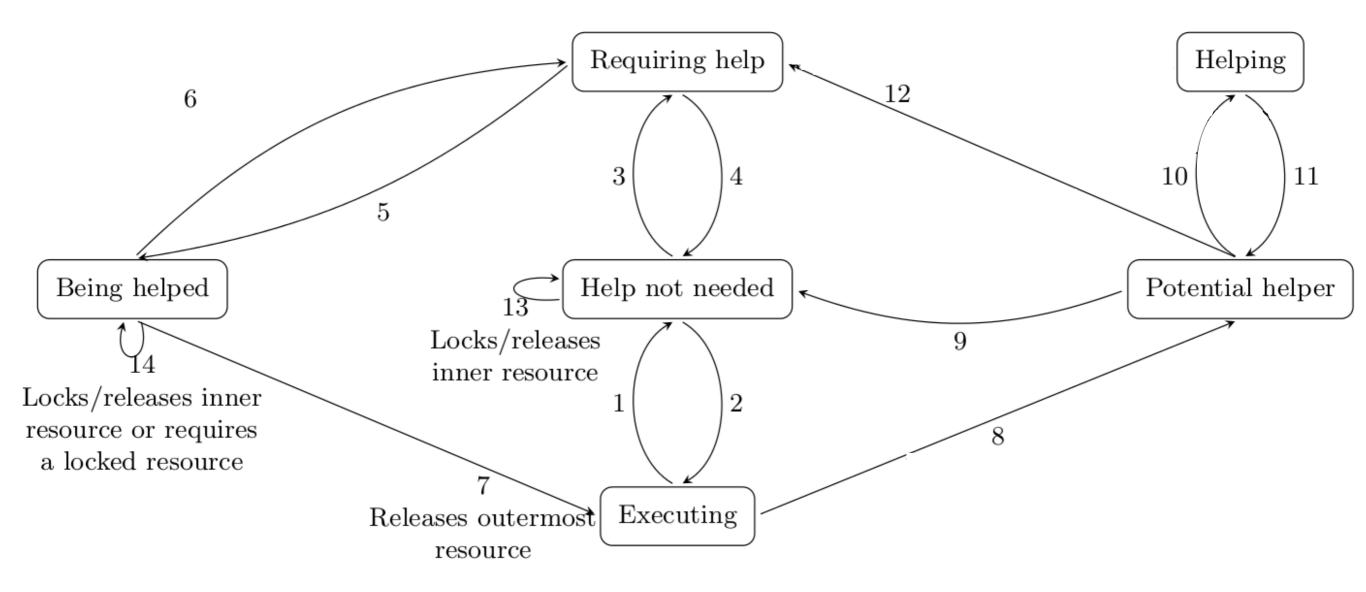
- Challenges
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 - Transitive blocking

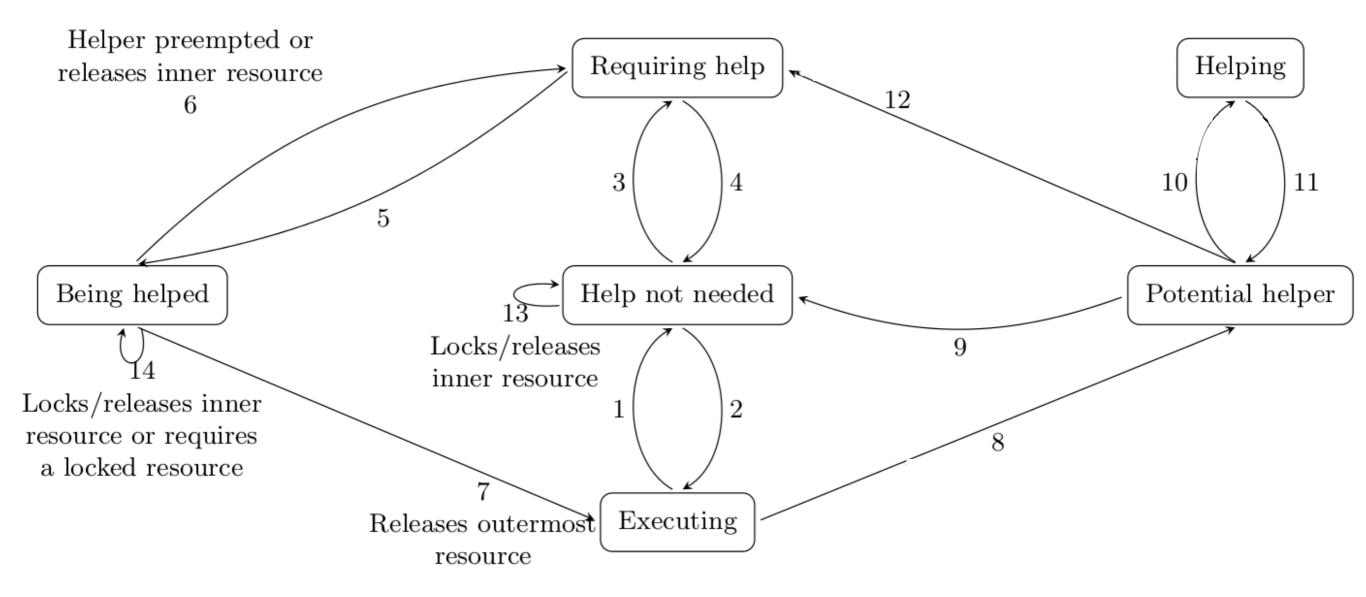
- Challenges
 - Global and Local resources
 - Bound concurrency
 - Transitive blocking
 - Local blocking

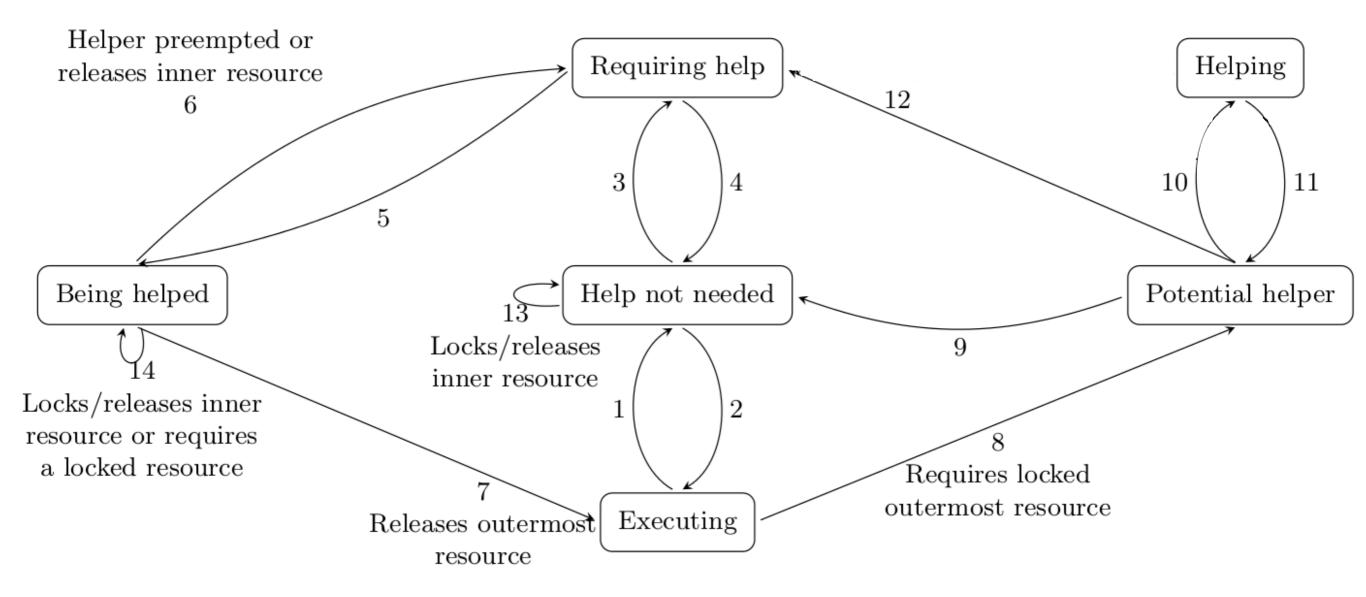


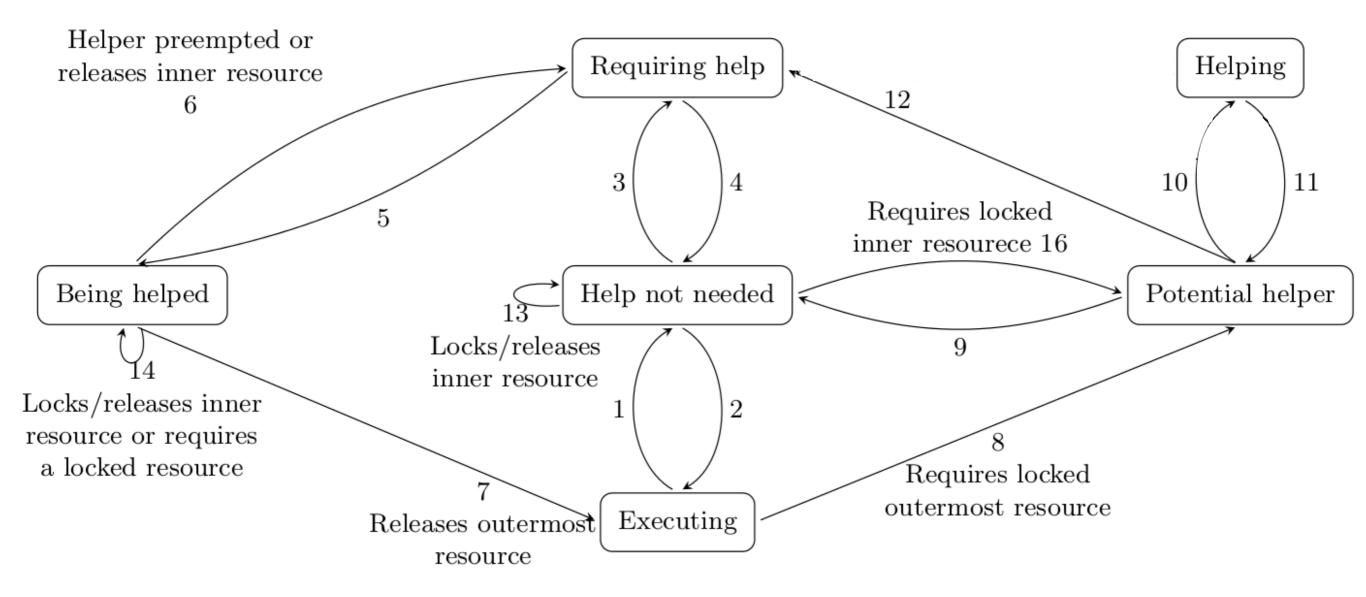


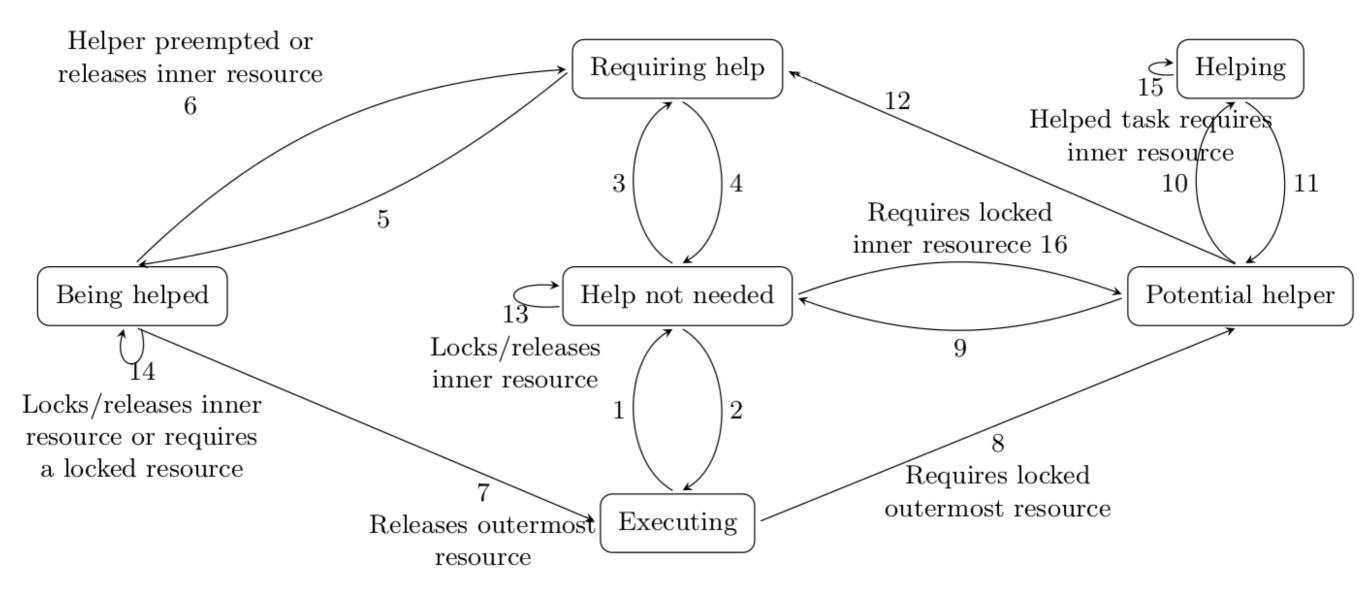


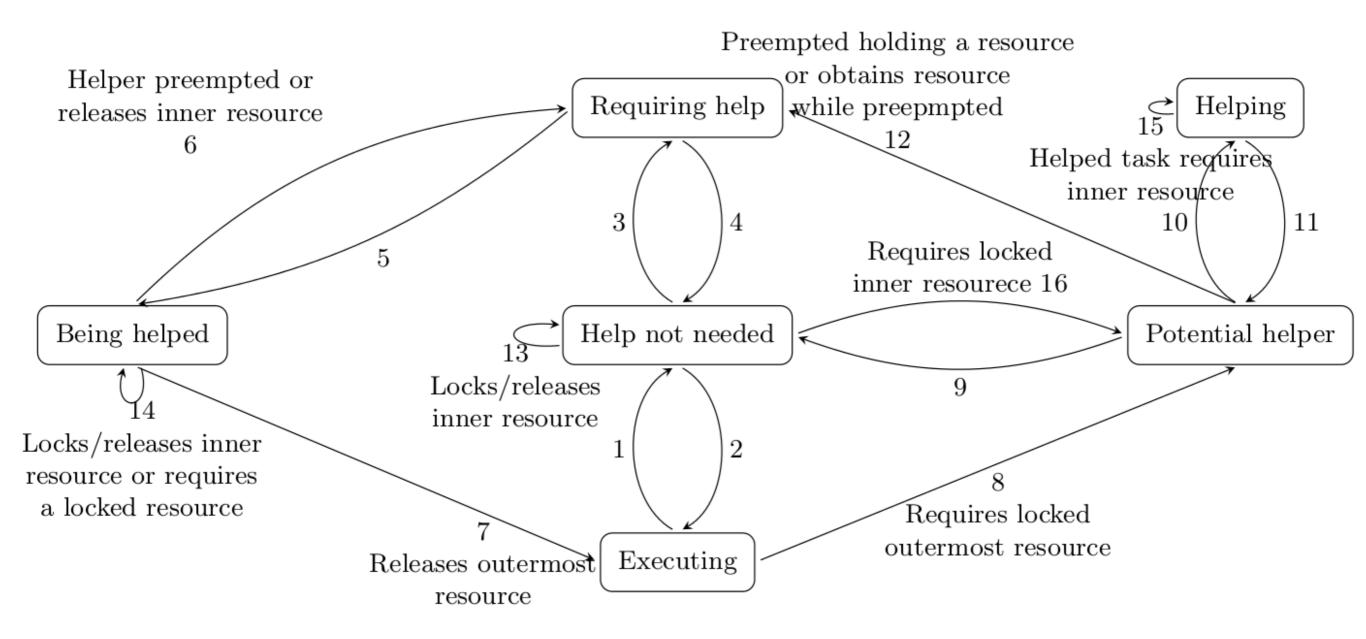












$$e^{j} = (|V(r^{j})| + |map(G(r^{j}))|) * (c^{j} + \sum_{r^{k} \in \mathbf{U}(r^{j})} n_{j}^{k} e^{k})$$

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Access request sources:

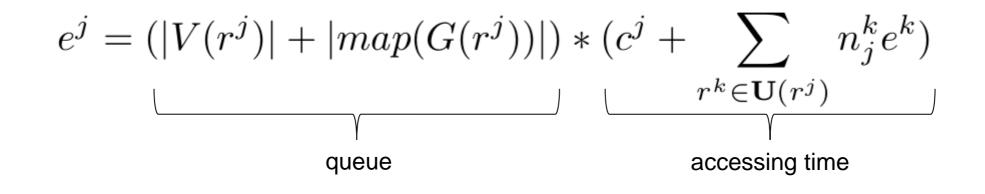
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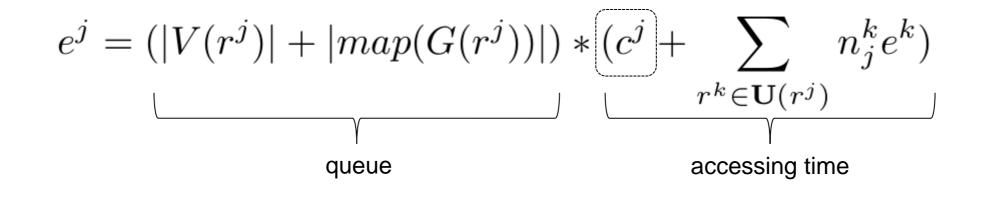
- Access request sources:
 - Tasks request access from inside another resource

$$e^{j} = (|V(r^{j})| + \underbrace{|map(G(r^{j}))|)}_{\text{queue}} * (c^{j} + \sum_{r^{k} \in \mathbf{U}(r^{j})} n_{j}^{k} e^{k})$$

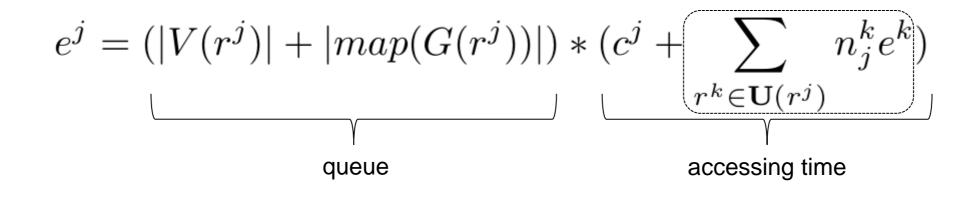
- Access request sources:
 - Tasks request access from inside another resource
 - Tasks not holding a resource



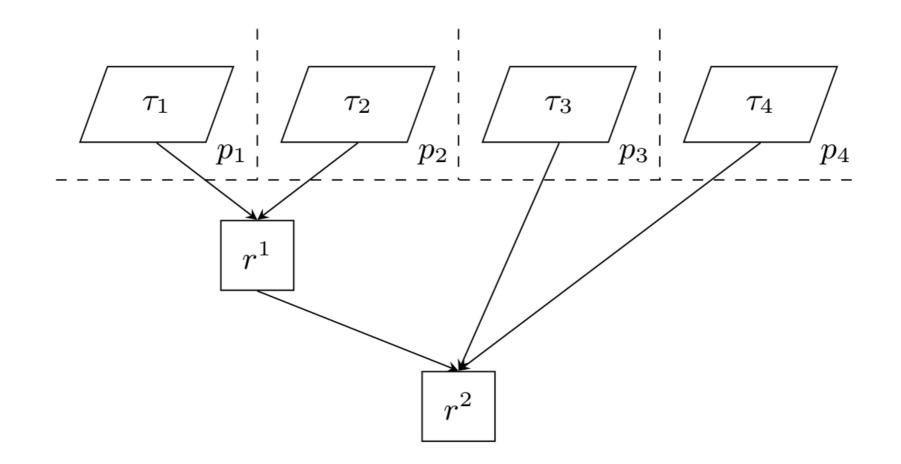
- Access request sources:
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- Access time:

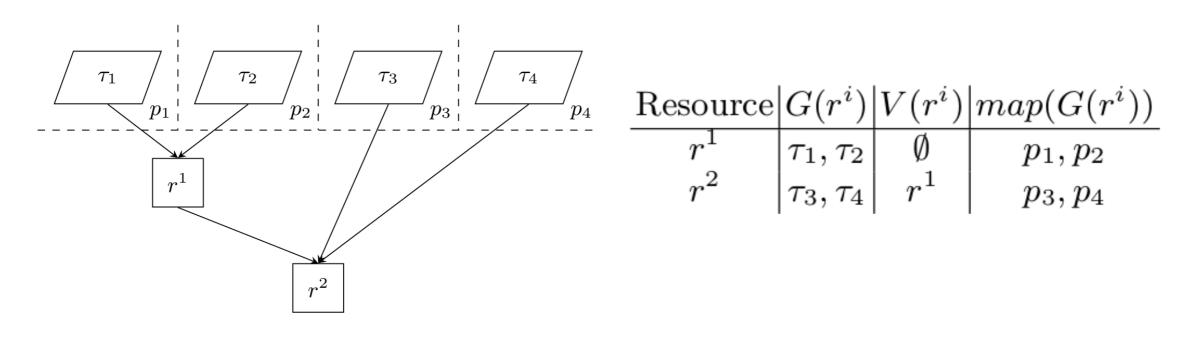


- Access request sources:
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- Access time:
 - Resource execution time

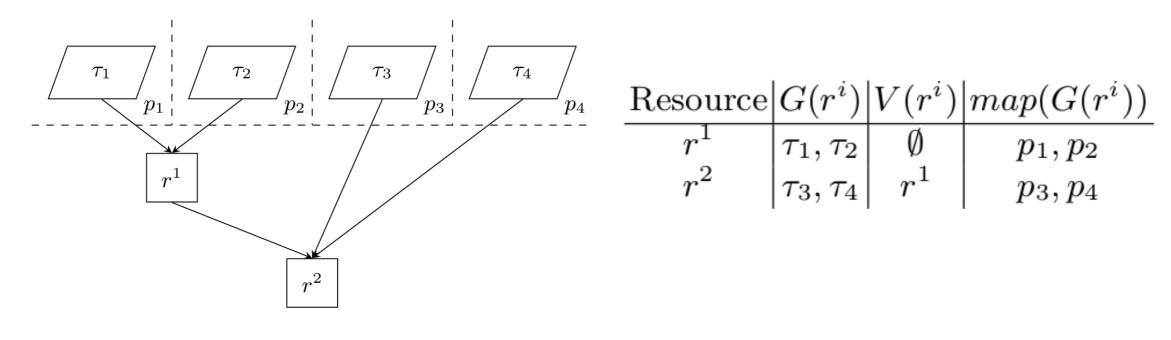


- Access request sources:
 - Tasks request access from inside another resource
 - Tasks not holding a resource
- Access time:
 - Resource execution time
 - Inner resources access costs



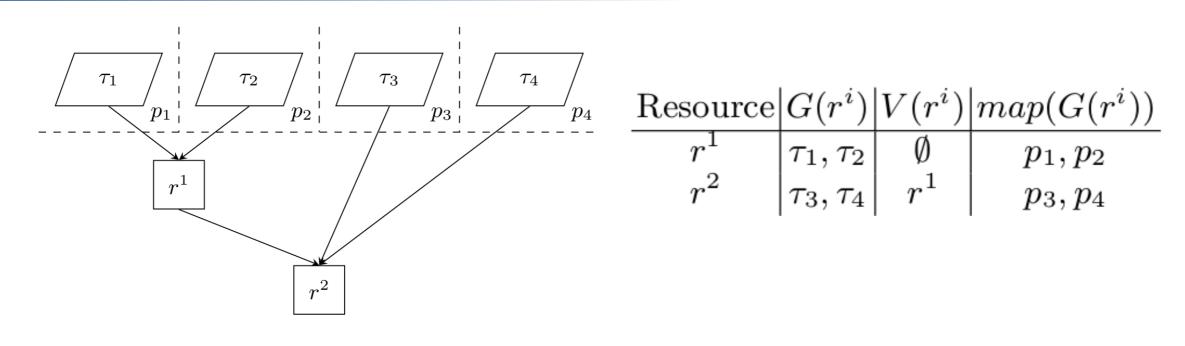


$$e^{j} = (|V(r^{j})| + |map(G(r^{j}))|) * (c^{j} + \sum_{r^{k} \in \mathbf{U}(r^{j})} n_{j}^{k} e^{k})$$



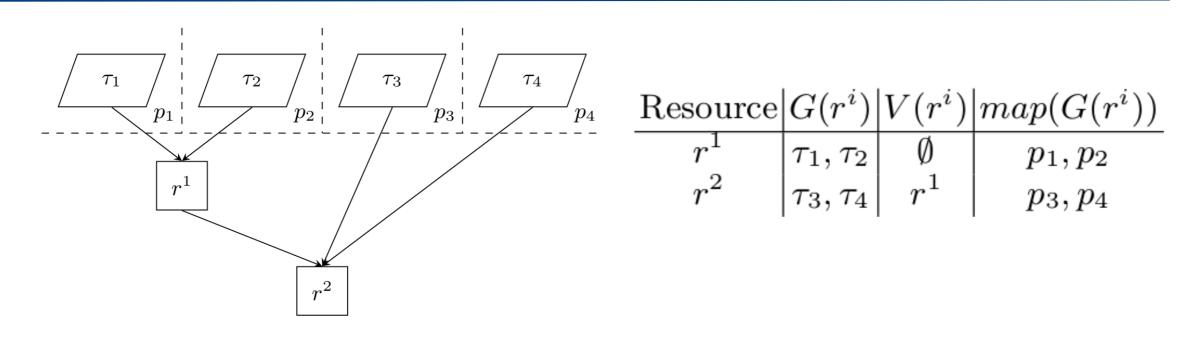
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$$e^{2} =$$



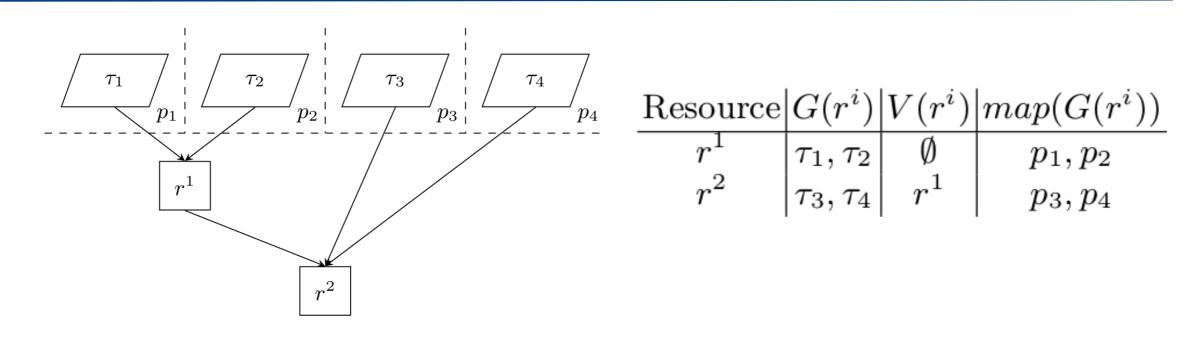
$$e^{j} = (|V(r^{j})| + |map(G(r^{j}))|) * (c^{j} + \sum_{r^{k} \in \mathbf{U}(r^{j})} n_{j}^{k} e^{k})$$

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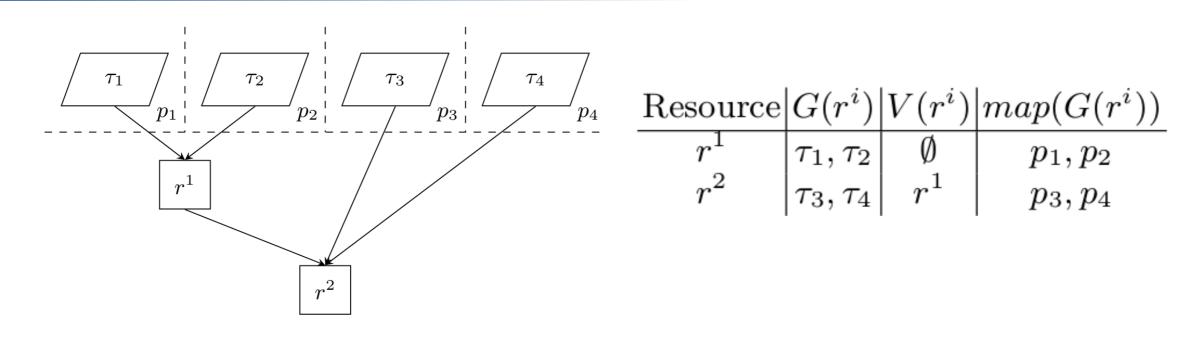
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$$e^2 = (1+2) *$$



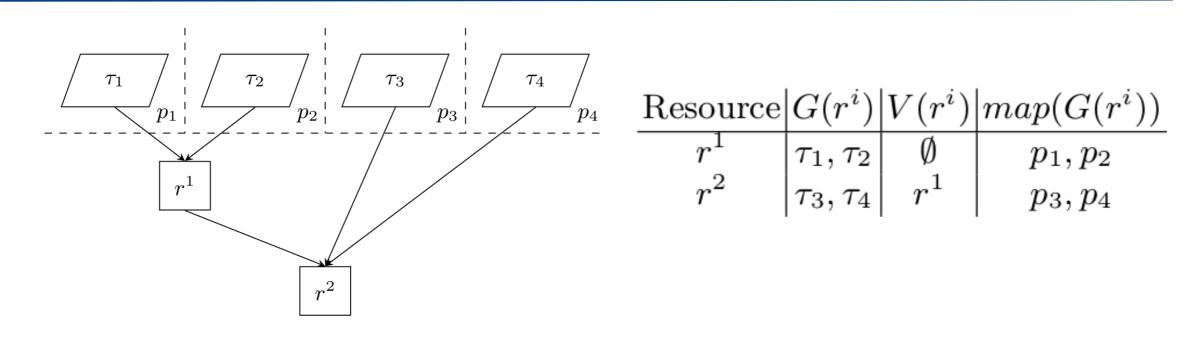
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$$e^2 = (1+2) * (c^2) =$$



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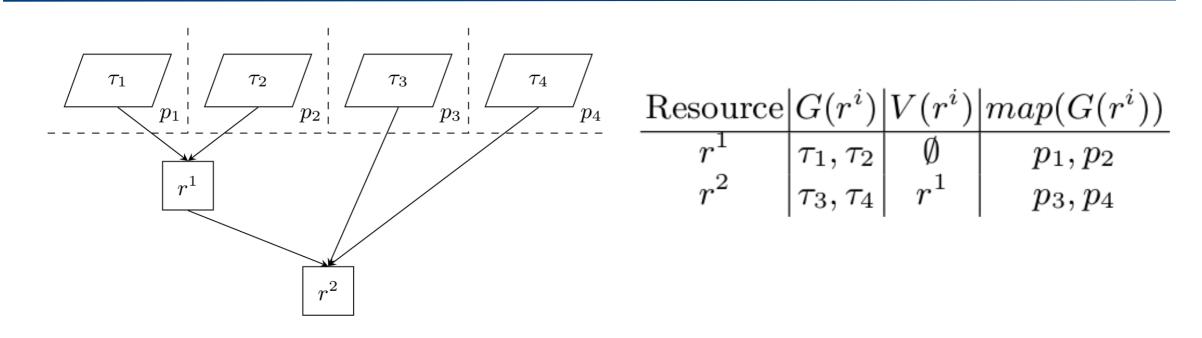
$$e^2 = (1+2) * (c^2) = 3c^2$$



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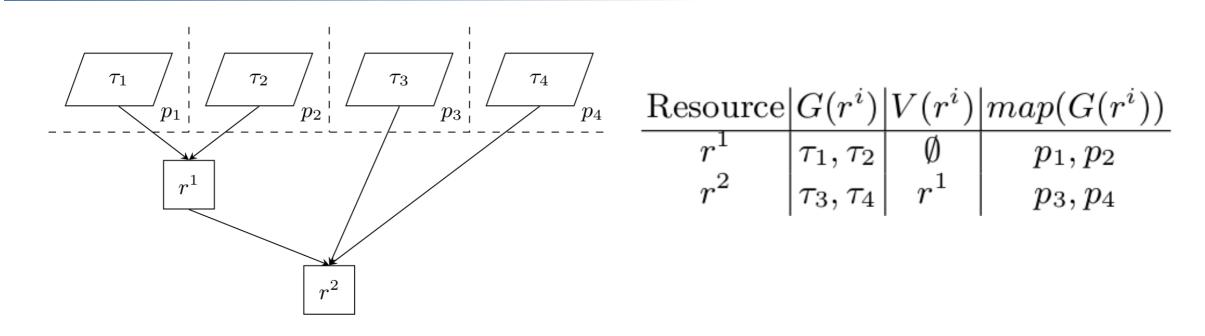
 $e^1 =$



$$e^{j} = (|V(r^{j})| + |map(G(r^{j}))|) * (c^{j} + \sum_{r^{k} \in \mathbf{U}(r^{j})} n_{j}^{k} e^{k})$$

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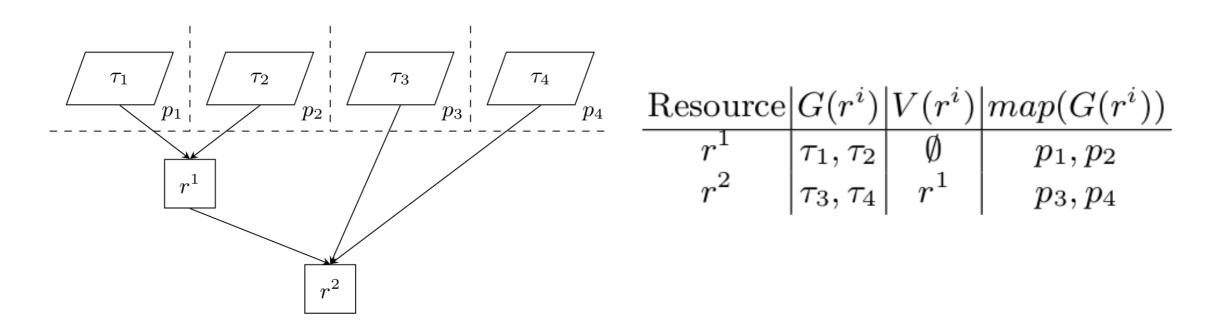
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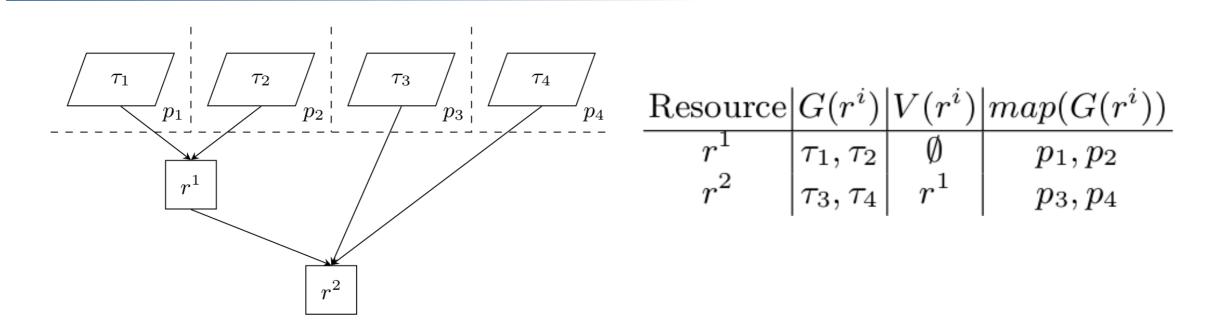
$$e^1 = (0+2) *$$



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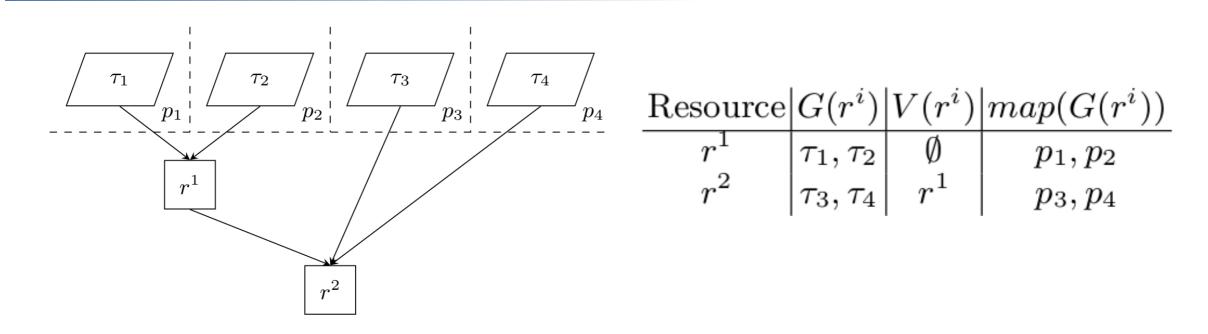
$$e^1 = (0+2) * (c^1 + c^2) + (c^2 + c^2) +$$



$$e^{j} = (|V(r^{j})| + |map(G(r^{j}))|) * (c^{j} + \sum_{r^{k} \in \mathbf{U}(r^{j})} n_{j}^{k} e^{k})$$

$$e^2 = (1+2) * (c^2) = 3c^2$$

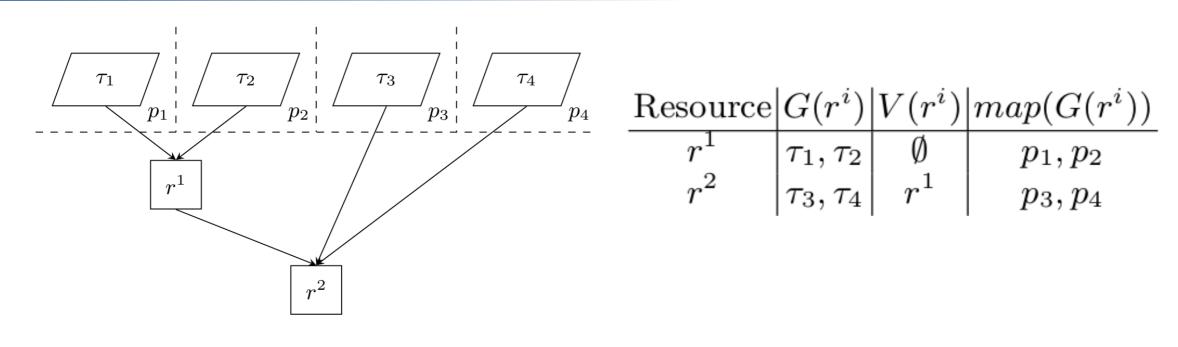
$$e^1 = (0+2) * (c^1 + e^2) =$$



 $e^{j} = (|V(r^{j})| + |map(G(r^{j}))|) * (c^{j} + \sum_{r^{k} \in \mathbf{U}(r^{j})} n_{j}^{k} e^{k})$

$$e^2 = (1+2) * (c^2) = 3c^2$$

$$e^1 = (0+2) * (c^1 + e^2) = 2(c^1 + e^2) =$$



$$e^{j} = (|V(r^{j})| + |map(G(r^{j}))|) * (c^{j} + \sum_{r^{k} \in \mathbf{U}(r^{j})} n_{j}^{k} e^{k})$$

$$e^2 = (1+2) * (c^2) = 3c^2$$

$$e^{1} = (0+2) * (c^{1} + e^{2}) = 2(c^{1} + e^{2}) = 2(c^{1} + 3c^{2})$$

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- Helped task executes at helper active priority
 - Can be different to the held resource local priority
- Tasks only update their priority on their own processor
- Transitive helping mechanism
 - If t3 is waiting for t2, which waiting for t1, if the later are locally preempted, t3 helps t1
- Tasks can help migrated local tasks
- Migrated tasks remain notionally dispatchable at host

MrsP – Conclusions

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Provided full nested resource model to MrsP

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- Provided full nested resource model to MrsP
- Updated timing analysis

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 - Nested shared resource access cost bounded

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 - Fine grained analysis





Supporting nested resources in MrsP

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