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## OpenMP tasking model for Ada: safety and correctness

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Vienna (Austria) June 12-16, 2017

#### Parallel heterogeneous embedded architectures

- 1. Exploit its performance capabilities
- 2. Facilitate programmability
- 3. Ensure portability



#### **Parallel Programming Models**



NVIDIA Tegra X1: 4-core A57 and A53, GPU (automotive)



Kalray MPPA: four 4-core K1, 256-core fabric (avionics)



TI Keystone II: 4-core A15, 8-core DSP (industrial)



# Parallel Programming Models

- Provides a level of abstraction to express parallelism while hiding processor complexities
  - Defines parallel regions and synchronization mechanisms
  - Couples host processor with accelerators devices (e.g., many-cores and DSP fabrics, GPUs, FPGAs)
- Mandatory as the number computing resources integrated increases

Host	
	Accelerator
	device(s)



Generic parallel heterogeneous architecture

# Ada and Parallel Programming Models

• There is a necessity to support parallelism in Ada capable of exploiting parallel heterogeneous embedded architectures

•			
	Define an Ada's parallel model <sup>1</sup>	Adopt an existing parallel model	
Pros	<ul> <li>Full control of the model</li> <li>Incorporate safety issues in the model</li> </ul>	<ul> <li>Very mature models</li> <li>Portability</li> <li>Develop "only" the Ada and parallel run-time connection</li> </ul>	
Cons	<ul> <li>Develop the complete parallel framework</li> <li>Less portability</li> </ul>	<ul> <li>No safety properties</li> </ul>	

<sup>1</sup> Michell, Moore, Pinho, *Tasklets – a fine grained parallelism for Ada on multicores*, in Ada-Europe 2013 Pinho, Moore, Michell, Taft, *Real-Time Fine-Grained Parallelism in Ada*, in ACM SIGAda Ada Letters 2015



## Parallel Programming Model Challenges

- Productivity perspective (performance, programmability, portability)
  - Shared and distributed memory
  - Fine-grain task- and data-based parallelism
  - Heterogeneity
  - Load balancing
  - Efficient synchronization methods



- Safety perspective
  - Parallel programming is complex and error prone, compromising correctness and so safety
  - Compiler and run-time techniques to detect errors must be provided







- Mature language constantly reviewed and augmented (last release Nov 2015)
- Performance and efficiency
  - Tantamount to other models (e.g. TBB, CUDA, OpenCL and MPI)
  - Support for fine-grain data- and task-parallelism
  - Features an advanced accelerator model for heterogeneous computing
- Portability
  - Supported by many chip and compiler vendors (Intel, IBM, ARM, TI, Kalray, Gaisler)
- Programmability
  - Currently available for C, C++ and Fortran (#pragma omp)
  - Allows incremental parallelization
  - Can be easily compiled sequentially (easing debugging)



## **OpenMP** execution model

- Fork-join parallel model of execution
- Task-centric model

It is important not to confuse *OpenMP tasks, Ada tasks* and *Ada tasklets* are not the same thing

- Ada tasks are meant to exploit concurrency
- OpenMP tasks and Ada tasklets are meant to exploit fine-grain parallelism





# **OpenMP** memory model

#### **Relaxed-consistency** memory model

Variables visibility defined by the programmer: shared, private, firstprivate



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# **OpenMP** and Safety

- Our vision is that OpenMP enables to guarantee safety requirements in terms of
  - Time predictability



- Reasoning about the timing behaviour of the parallel execution
- Safety and correctness
  - Ensuring that the correct operation in response to its inputs
  - Support reliability and resiliency mechanisms

Compiler analysis techniques for checking correctness **Error handling** methodologies to be added in the specification



## Ada and OpenMP

### **Our proposal**

- 1. Extend OpenMP to support Ada
- 2. Extend Ada to support OpenMP (e.g., including a new pragma OMP)
- 3. Add compiler and runtime mechanisms to ensure **correctness**

#### Example: Fibonacci computation

```
function Fibonacci (N: Integer) return Integer is
begin
  if N < 2 then
    return N;
  pragma OMP (parallel, shared=>X,Y,
                        firstprivate=>N);
  pragma OMP (single, nowait);
  begin
    pragma OMP (task, shared=>X,
                       firstprivate=>N);
      X := Fibonacci(N - 2);
    pragma OMP (task, shared=>Y,
                       firstprivate=>N)
      Y := Fibonacci(N - 2);
  end
  return X + Y;
end Fibonacci;
```



## Fine-grained parallelism in Ada

	Blocks	Loops	Reductions	Tasks
Tasklet Model <sup>1</sup>	<pre>parallel   seq_of_stat_1   and    seq_of_stat_2   end parallel;</pre>	<pre>for i in parallel lbub loop   seq_of_stat end loop;</pre>	<pre>type t is new array   (parallel &lt;&gt;) of Float   with Reducer =&gt; "+", Identity =&gt; 0.0; Par_Sum : t := (others =&gt; 0.0); begin   for I in parallel Arr'Range loop     seq_of_stat   end loop;   Sum := Par_Sum(&lt;&gt;)'Reduced;</pre>	-
OpenMP	<pre>pragma OMP (parallel); pragma OMP (single); begin pragma OMP (task); seq_of_stat_1 pragma OMP (task); seq_of_stat_2 end</pre>	<pre>pragma OMP (parallel); pragma OMP (taskloop); for i in range lbub loop seq_of_stat end loop;</pre>	<pre>pragma OMP (parallel); pragma OMP (taskloop, reduction=&gt;+,TOTAL); begin for i in range 0MAX_I loop seq_of_stat end loop; end</pre>	<pre>pragma OMP (parallel); pragma OMP (single); begin if cond then pragma OMP (task); seq_of_stat_1 else pragma OMP (task); seq_of_stat_2 end if; end</pre>

<sup>1</sup> Michell, Moore, Pinho, *Tasklets – a fine grained parallelism for Ada on multicores*, in Ada-Europe 2013 Pinho, Moore, Michell, Taft, *Real-Time Fine-Grained Parallelism in Ada*, in ACM SIGAda Ada Letters 2015

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- 🗸 Non-determinism
- 2. Race conditions
- 3. Deadlocks
- 4. Fault tolerance



- Actions not defined by the specification, e.g.,
  - Directives/Clauses receives arguments out of range
    - E.g., num\_threads (N)
  - Where and how some expressions have to be executed is not defined
    - E.g., the order in which the values of a reductions are combined
  - Compilers and runtimes are not forced to check the conformity of a program
    - E.g., the storage location specified in task dependencies must be identical or disjoint

<sup>1</sup>A functional safety OpenMP for critical real-time embedded systems

In the 13th International Workshop on OpenMP (IWOMP), New York (USA), September 18-19, 2017 Sara Royuela, Alejandro Duran, Maria A. Serrano, Eduardo Quiñones, Xavier Martorell



- Non-determinism
- Incorrect data scoping definition
- **Race conditions** 2.
- Deadlocks 3.
- Fault tolerance 4.

Incorrect usage of synchronization mechanism





- ✓ Non-determinism
- Race conditions
- 3. Deadlocks
- 4. Fault tolerance

There exist compilation techniques capable of identifying (and solving) race conditions<sup>1,2</sup>

• Incorrect usage of synchronization mechanism

Incorrect data scoping definition

```
int a = 1, b = 1, res;
int foo() {
  #pragma omp parallel shared(res) firstprivate(a,b)
  #pragma omp single
    int x, y;
    #pragma omp task shared(x) firstprivate(a)
     x = a*a;
    #pragma omp task shared(y) firstprivate(b)
     y = b*b;
    #pragma omp taskwait
    res = x + y;
  return res;
```

<sup>1</sup> Royuela, Duran, Liao, Quinlan, *Auto-scoping for OpenMP tasks*, in IWOMP 2012 <sup>2</sup> Lin, *Static nonconcurrency analysis of openmp programs*, in IWOMP 2008



- Non-determinism
- Race conditions
- ✓ Deadlocks
- 4. Fault tolerance



OpenMP synchronization mechanisms might result in

- Possible solutions to avoid deadlocks
  - Check that programs are OpenMP conformant
  - Adapt already existing compiler methods to OpenMP<sup>1</sup>
  - Avoid OpenMP techniques in favor of Ada high-level concurrency mechanisms (e.g., protected objects)





- ✓ Non-determinism
   ✓ Race conditions
   ✓ Deadlocks
   ✓ Fault tolerance
- One major problem of OpenMP in safety environments is the lack of resiliency mechanisms
- Attemps to add error-handling mechanisms to the standard already exist<sup>1</sup>
  - Some proposals have already been adopted (cancellation constructs)

<sup>1</sup> Duran et. al., *A proposal for error handling in OpenMP* Wong et. al., *Towards an error model for OpenMP* Fan et. al., *Exception handling with OpenMP in object oriented languages* 





- There is a necessity to extend Ada with fine-grained parallelism to efficiently support parallel heterogeneous computing
- Our proposal: To adopt OpenMP as a parallel programming model for Ada
  - Very mature parallel programming model (20 years)
  - Performance, programmability and portability without jeopardizing safety
  - Parallel programming challenges regarding safety can be addressed





- This can be complementary and compatible with the parallel Ada model
  - OpenMP tasks and Ada tasklets are similar
  - The interplay between Ada and OpenMP runtimes must be analyzed (e.g. OpenMP could be the runtime common to both)
  - The Ada community **can influence** the OpenMP standard to address the challenges that impacts on safety



# A proposal to extend OpenMP

#### A functional safety OpenMP for critical real-time embedded systems,

#### To be presented in the 13th International Workshop on OpenMP, celebrated in New York in September 18-19, 2017

Sara Royuela, Alejandro Duran, Maria A. Serrano, Eduardo Quiñones, Xavier Martorell

#### • Comments from reviewers

- "[..] the proposed extensions are a good step toward making the use of OpenMP in safety environments practical, and appear to provide real value [..]"
- "[..] Even if OpenMP didn't care about embedded systems this analysis seems useful to help elucidate some of the issues inherent in the OpenMP specification [..]"
- "[..] it is an interesting challenge for modification on current OpenMP [..] OpenMP ARB may consider this proposal [..]"



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## OpenMP tasking model for Ada: safety and correctness

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Туре	Language	V Strengths	🗙 Weaknesses
Hardware Centric	Intel <sup>®</sup> TBB	- Highly tunable - High-level (task concept)	<ul> <li>Portability</li> <li>Mapping thread/core not part of the model</li> </ul>
	NVIDIA® CUDA	<ul> <li>Highly tunable</li> <li>Wrappers for many languages</li> </ul>	<ul> <li>Low level (explicit data management)</li> <li>Restricted to NVIDIA GPUs</li> </ul>
Application Centric	OpenCL	<ul> <li>Automatic vectorization</li> <li>Executes in host and accelerator</li> </ul>	<ul> <li>Low level (explicit data management)</li> <li>Full rewriting</li> </ul>
	Pthreads	<ul> <li>Full execution control (thread concept)</li> <li>Dynamic creation/destruction of threads</li> </ul>	<ul> <li>Low level (reductions, work distribution, synchronization, etc. by hand)</li> </ul>
Parallelism Centric	OpenMP	<ul> <li>High-level (task and data-flow concept)</li> <li>Portable</li> <li>Exploits parallelism at host and device</li> </ul>	- No safety concepts



#### **Race conditions: general techniques**

Dynamic	Analyze specific executions and possibly deliver false negatives There are algorithms capable of detecting at least one race when races are present
Static	<ul> <li>Analyze all possibilities (NP-hard) and possibly deliver false positives</li> <li>Sound solutions exist for specific subsets of OpenMP</li> <li>Fixed number of threads</li> <li>Using affine constructs</li> <li>Also solutions for detecting non-concurrency</li> </ul>
Hybrid	Combination of static and dynamic tool for more accurate results

